

Bandwidth of graphs resulting from the edge clique covering problem

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Abstract

Let n, k, b be integers with $1 \leq k-1 \leq b \leq n$ and let $G_{n,k,b}$ be the graph whose vertices are the k -element subsets X of $\{0, \dots, n\}$ with $\max(X) - \min(X) \leq b$ and where two such vertices X, Y are joined by an edge if $\max(X \cup Y) - \min(X \cup Y) \leq b$. These graphs are generated by applying a transformation to maximal k -uniform hypergraphs of bandwidth b that is used to reduce the (weak) edge clique covering problem to a vertex clique covering problem. The bandwidth of $G_{n,k,b}$ is thus the largest possible bandwidth of any transformed k -uniform hypergraph of bandwidth b . For $b \geq \frac{n+k-1}{2}$, the exact bandwidth of these graphs is determined. For $b < \frac{n+k-1}{2}$, the bandwidth is asymptotically determined in the case of $b = o(n)$ and in the case of b growing linearly in n with a factor $\beta \in (0, 0.5]$, where for one case only bounds could be found. It is conjectured that the upper bound of this open case is the right asymptotic value.

1 Introduction

The bandwidth problem for graphs is to find a labelling of the vertices with different integers, such that the maximum absolute value of the difference of the labels of two adjacent vertices is minimal. There are many applications such as efficient storage of sparsely populated symmetric matrices, which arise e.g. from discretization of partial differential equations, cf. [20]. Several other applications, including the placement problem for modules of a VLSI design, the binary constraint satisfaction problem and the minimization of effects of noise in the multichannel communication of data are discussed e.g. in [5, 3, 8]. The bandwidth problem was shown to be NP-hard [18] and even an approximation with a ratio better than 2 is NP-hard [10], so several heuristics such as the Cuthill-McKee-Algorithm [9] or some similar approaches, cf. [12], are very popular in applications. However, for some graph classes the exact bandwidth is known. These include the path, the cycle, the complete graph, the complete bipartite graph [6], the hypercube [13], the grid graph [7], special Hamming

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graphs [16] and several others, cf. [5]. However, there are still many graphs, where the exact bandwidth is unknown, such as the general Hamming graphs, cf. [15, 2]. In this paper, we consider graphs $G_{n,k,b}$, $1 \leq k-1 \leq b \leq n$, whose vertices are those k -element subsets of $\{0, \dots, n\}$, for which the difference of the maximum and the minimum is at most b . There is an edge between two vertices, if the difference of the maximum and the minimum of the union of the corresponding sets is at most b . We start by introducing the necessary notation and a motivation in Section 2 and study some basic properties in Section 3. Based on that, we determine the exact bandwidth for these graphs in the case of $b \geq \frac{n+k-1}{2}$ in Section 4. In Section 5, we present some asymptotic results for $n \rightarrow \infty$ in the case of $b = o(n)$. The results of Sections 4 and 5 are summarized by the following theorem:

Theorem 1. *Let k be a fixed positive integer and $1 \leq k-1 \leq b \leq n$.*

a) *If $b \geq \frac{n+k-1}{2}$ then*

$$B(G_{n,k,b}) = \left\lceil \frac{(n+1)\binom{b}{k-1} - (k-1)\binom{b+1}{k} + \binom{2b-n+1}{k} - 2}{2} \right\rceil.$$

b) *If $b = o(n^{\frac{1}{k+1}})$ then for sufficiently large n*

$$B(G_{n,k,b}) = k\binom{b}{k}.$$

If $b = o(n)$ then

$$B(G_{n,k,b}) \sim k\binom{b}{k} \text{ as } n \rightarrow \infty.$$

Sections 6 to 9 discuss the case $b \sim \beta n$ with $\beta \in (0, 0.5]$. The main result is given by the next theorem:

Theorem 2. *Let $k \geq 2$ be a fixed positive integer, but $n \rightarrow \infty$. Let $b \sim \beta n$ and let $1 = q\beta + r$, where $q \geq 2$ is a positive integer and $0 \leq r < \beta$. Let*

$$\begin{aligned} c_1(\beta, k) &= \frac{\beta^k}{k!} \left(k - \frac{k-1}{q} \right), \\ c_2(\beta, k) &= \frac{\beta^{k-1}}{(q+1)k!} (k - (k-1)\beta), \\ c_3(\beta, k) &= \frac{(\beta-r)^k}{(q+1)k!} q^{k-1}. \end{aligned}$$

a) *If $r \leq \frac{q-1}{q^2+q-1}$ then $B(G_{n,k,b}) \sim c_1(\beta, k)n^k$.*

b) *If $r > \frac{q-1}{q^2+q-1}$ then $\max\{c_1(\beta, k), c_2(\beta, k) + \frac{1}{q^{k-1}}c_3(\beta, k)\}n^k \lesssim B(G_{n,k,b}) \lesssim (c_2(\beta, k) + c_3(\beta, k))n^k$.*

The part b) gives only bounds instead of an exact asymptotic value. We strongly conjecture that the RHS bound is the right value. The bounds are not too far away from each other because

$$\frac{c_2(\beta, k)}{c_3(\beta, k)} = \left(\frac{\beta}{q\beta - qr} \right)^k \frac{q}{\beta} (k - (k-1)\beta) \geq \left(\frac{\beta}{q\beta - qr} \right)^k \frac{k(q-1) + 1}{\beta} \geq 6$$

since $\beta > q\beta - qr$ iff $r > \frac{q-1}{q^2+q-1}$ and $k \geq 2, q \geq 2, \beta \leq \frac{1}{2}$.

Let $U = \{\beta \in (0, 0.5] : r > \frac{q-1}{q^2+q-1}\}$ be the set of numbers β for which part b) applies and thus the exact asymptotic value is still unknown. Note that $r > \frac{q-1}{q^2+q-1}$ iff $\frac{1}{q+1} < \beta < \frac{q}{q^2+q-1}$. Thus the Lebesgue measure of U is equal to $\sum_{q=2}^{\infty} (\frac{q}{q^2+q-1} - \frac{1}{q+1}) = 0.119\dots$, i.e., for the “majority” of numbers $\beta \in (0, 0.5]$ the exact value is known.

The proof of Theorem 2 is based on a reduction to a continuous problem on the unit square $[0, 1]^2$. Riemann integrals and elementary geometric arguments suffice. The embedding into a more difficult continuous problem on the unit cube was used by Harper [14] to obtain bounds for the bandwidth of Hamming graphs. Also for the edge-bandwidth of multidimensional grids and Hamming graphs (the bandwidth of the line graph of these graphs) Harper’s reduction to the unit cube was applied in [1]. Asymptotic bounds for the bandwidth of the d -ary de Bruijn graph were obtained in [19] by an approach based on the use of a continuous domain.

2 Notation and motivation

Let $[n] = \{1, \dots, n\}$ and $[i, j] = \{i, i+1, \dots, j-1, j\}$ with $i, j \in \mathbb{Z}, i \leq j$. In particular, $[0, n] = \{0, 1, \dots, n\}$. For a graph $G = (V, E)$ with $|V| = n$ vertices, a *proper numbering* of G is a bijection $f : V \rightarrow [n]$. For two vertices $u, v \in V$, we call $d_f(u, v) = |f(u) - f(v)|$ the f -distance of u and v . Let f be a proper numbering of a graph G . The *bandwidth of f* , denoted $B_f(G)$, is given by $B_f(G) = \max\{|f(u) - f(v)| : \{u, v\} \in E\}$, i.e., the maximal f -distance. The *bandwidth of G* is defined by $B(G) = \min\{B_f(G) : f \text{ is a proper numbering of } G\}$. A *bandwidth numbering* of G is a numbering f such that $B(G) = B_f(G)$. This definition can be easily generalized to hypergraphs $H = (V, E)$. There we have $B_f(H) = \max\{|f(u) - f(v)| : \exists e \in E \text{ with } u, v \in e\}$.

Now we formally define the subject of our study. Let k and b be positive integers with $b \geq k - 1$. For $A \subseteq [0, n]$ let $\underline{A} = \min(A)$ and $\overline{A} = \max(A)$. Further let $\binom{[0, n]}{k} = \{X \subseteq [0, n] : |X| = k\}$. Then $G_{n, k, b}$ is the graph with vertex set

$$V_{n, k, b} = \left\{ X \in \binom{[0, n]}{k} : \overline{X} - \underline{X} \leq b \right\}$$

and edge set

$$E_{n, k, b} = \left\{ \{X, Y\} \in V_{n, k, b} : \overline{X \cup Y} - \underline{X \cup Y} \leq b \right\}.$$

These graphs arise in the following context: Let $G = (V, E)$ be a graph. A *clique* is a subset of V that induces a complete subgraph of G . Consider the following transformation, which was used in the NP-completeness proof of the edge clique covering problem in [17] by reducing it to the vertex clique covering problem. Let $V = [n]$. Then $\tilde{G} = (\tilde{V}, \tilde{E})$ is constructed by setting $\tilde{V} = E$ and $\{\{i, j\}, \{i', j'\}\} \in \tilde{E}$, if there is a clique $C \subseteq V$ with $i, j, i', j' \in C$. Let $\bar{\chi}_v(G)$ (resp. $\bar{\chi}_e(G)$) be the *vertex clique covering number* (resp. *edge clique covering number*) of the graph $G = (V, E)$, i.e., the minimal number of cliques whose induced subgraphs cover all vertices (resp. edges) of G . It can be shown, that $\bar{\chi}_e(G) = \bar{\chi}_v(\tilde{G})$, which is the essential part of the NP-completeness proof for the edge clique covering problem, since the transformation can be done in polynomial time. Here we want to generalize this result for hypergraphs $H = (V, E)$. We consider the *2-section graph* of H , i.e., the graph $G_H = (V, E_H)$ on the vertex set of H where $\{u, v\} \in E_H$ if there is an edge of H containing u and v . A subset C of V is called a *weak clique* of H if C is a clique in G_H . A *weak edge clique covering* of H is a family \mathcal{C} of weak cliques of H such that for all $e \in E$ there is some $C \in \mathcal{C}$ with $e \subseteq C$. The *weak edge clique covering number* of H is the smallest size $\bar{\chi}_e(H)$ of a weak edge clique covering of H . It will turn out that the computation of $\bar{\chi}_e(H)$ can be simplified by the following transformation. The *weak edge clique graph* of H is the graph $\tilde{G}_H = (\tilde{V}, \tilde{E})$ where $\tilde{V} = E$ and two vertices of \tilde{G}_H , i.e., edges e, e' of H , are adjacent if there is a weak clique C of H containing e and e' as subsets. With these definitions we prove the following proposition.

Proposition 1. *Let $H = (V, E)$ be a hypergraph. Then*

$$\bar{\chi}_e(H) = \bar{\chi}_v(\tilde{G}_H).$$

Proof. Let \mathcal{C} be a weak edge clique covering of H of size $\bar{\chi}_e(H)$. For each $C \in \mathcal{C}$ let $\tilde{C} = \{e \in E : e \subseteq C\}$. Then \tilde{C} is a clique in \tilde{G}_H and $\tilde{\mathcal{C}} = \{\tilde{C} : C \in \mathcal{C}\}$ is a vertex clique covering of \tilde{G}_H . Consequently $\bar{\chi}_e(H) \geq \bar{\chi}_v(\tilde{G}_H)$. Now let $\tilde{\mathcal{C}}$ be a vertex clique covering of \tilde{G}_H of size $\bar{\chi}_v(\tilde{G}_H)$. For each $\tilde{C} \in \tilde{\mathcal{C}}$ let $C = \bigcup_{e \in \tilde{C}} e$. Then C is a weak clique in H . To verify this fact, we pick two arbitrary vertices $x, y \in C$ and show that they are adjacent in G_H . First we consider the case that there is an edge $e \in \tilde{C}$ with $x, y \in e$. Then, by construction of G_H , $\{x, y\} \in E_H$. The alternative is, that there are 2 edges $e, e' \in \tilde{C}$ with $x \in e$ and $y \in e'$. Since e and e' are adjacent in \tilde{G}_H there is a weak clique of H containing both edges, which implies the adjacency of x and y . Moreover, $\tilde{\mathcal{C}}$ is a weak edge clique covering of H . Consequently, $\bar{\chi}_v(\tilde{G}_H) \geq \bar{\chi}_e(H)$. \square

Thus, from an algorithmic point of view, it is enough to study the vertex clique covering problem. For bounded bandwidth, and more generally for bounded treewidth, there is a linear time dynamic programming algorithm for the solution [4]. In an application, which will be described below, we were lead to the weak edge covering problem on a hypergraph whose bandwidth is small (and thus, theoretically, considered as bounded). This implies the following question: Given a hypergraph H of bandwidth b , how large can be the bandwidth of the weak edge clique graph \tilde{G}_H of H ? Here we discuss only k -uniform hypergraphs

though many results can be simply generalized to hypergraphs whose edges have size at most k . For later computations, it is more suitable to take $[0, n]$ as the vertex set of H instead of $[1, n]$. If, without loss of generality, $f(i) = i$ is the bandwidth numbering of H then, obviously, \tilde{G}_H has maximal bandwidth if H contains all k -element subsets X of $[0, n]$ with $\overline{X} - \underline{X} \leq b$. In this case, \tilde{G}_H is exactly the graph $G_{n,k,b}$, which motivates the study of $G_{n,k,b}$.

We came to these questions in the study of multielectrode recordings of neuronal signals, so-called spikes. Such recordings are carried out on multielectrode arrays, which can be used in-vivo or in-vitro. The denser the electrodes are placed the more likely it is for the neurons to be simultaneously recorded at different electrodes. The resulting similarities in the recordings of the electrodes can provide useful information. In [11] we developed an algorithm to estimate the (unknown) neighborhood of a neuron, i.e., the set of electrodes which record the signals of this neuron. Such neighborhood information is also used as an additional tool in [21] for the so called *spike sorting*, which is an estimated assignment of the recorded signals to the neurons.

Fix a short time interval in which several electrodes record signals. We consider these electrodes as vertices of a graph, which we call *similarity-graph* for the fixed time interval. First we mention that some neurons may always spike simultaneously. We combine such a set of neurons to one (artificial) new neuron. It might be an accident that two electrodes record a signal at almost the same time, but the simultaneous recording can also be caused by the fact that one spiking neuron has contact to both electrodes. Thus we do not test only one short time interval but several such intervals. If there are sufficiently many simultaneous recordings of two (or k) fixed electrodes, one may expect that these recordings are indeed caused by only one neuron and thus we draw an edge (hyperedge) between the corresponding vertices in the similarity-graph.

By algorithmic reasons, it is easier to check only pairs of electrodes, see [11]. But, with some more effort, also k -element subsets of electrodes could be checked for similarities if k is small. This leads to edges and hyperedges of electrodes. If a spiking neuron has contact to an unknown set S of electrodes, all edges between any two vertices of S (all hyperedges of any k vertices of S) are drawn in the similarity-graph. Though these edges may also be caused by different simultaneously spiking neurons having contact in each case to two (or k) neurons, it is more likely that only one neuron is the source. Such a neuron yields the edges of a weak clique in the similarity-graph. Once the similarity-graph is constructed, it remains the question what is the basic cause for this graph. A reasonable answer is that as few as possible neurons yield the graph. Consequently, a minimum weak edge clique covering has to be determined. Because of the bounded length of the axons, only nearby electrodes, which are placed in form of a two-dimensional bounded grid (or some similar variants), may have contact to the same neuron. Hence the similarity-graph is a relatively sparse graph and edges are only drawn between electrodes which have a small Euclidean distance. Thus it is reasonable to expect that also this graph has a small bandwidth.

3 Some basic properties

Lemma 1. *Let X and Y be two distinct vertices of $G_{n,k,b}$. They are adjacent iff $\overline{X} - \underline{Y} \leq b$ and $\overline{Y} - \underline{X} \leq b$.*

Proof. Let X and Y be adjacent. Then $\overline{X} - \underline{Y} \leq \overline{X \cup Y} - \underline{X \cup Y} \leq b$ and, analogously, $\overline{Y} - \underline{X} \leq b$.

Now let $\overline{X} - \underline{Y} \leq b$ and $\overline{Y} - \underline{X} \leq b$. Then $\overline{X \cup Y} - \underline{X \cup Y} = \max\{\overline{X}, \overline{Y}\} - \min\{\underline{X}, \underline{Y}\} = \max\{\overline{X} - \underline{X}, \overline{X} - \underline{Y}, \overline{Y} - \underline{X}, \overline{Y} - \underline{Y}\} \leq b$. \square

Note that $[i, i+k-1] \in V_{n,k,b}$ iff $0 \leq i \leq n-(k-1)$ and that for $0 \leq i < j \leq n-(k-1)$ the vertices $[i, i+k-1]$ and $[j, j+k-1]$ are adjacent iff $j \leq i+b-(k-1)$.

Lemma 2. *Let $0 \leq i < j \leq n-(k-1)$. The vertices $[i, i+k-1]$ and $[j, j+k-1]$ have distance $\lceil \frac{j-i}{b-k+1} \rceil$ in $G_{n,k,b}$.*

Proof. Let $j-i = q(b-k+1) + r$ where q is an integer and $1 \leq r \leq b-k+1$. Then $\lceil \frac{j-i}{b-k+1} \rceil = q+1$. Obviously, the vertices $[i, i+(k-1)], [i+b-(k-1), i+b], [i+2b-2(k-1), i+2b-(k-1)], \dots, [i+qb-q(k-1), i+qb-(q-1)(k-1)], [j, j+(k-1)]$ form a path in $G_{n,k,b}$ of length $q+1$. Thus the distance is at most $q+1$.

If the vertices $X_0 = [i, i+(k-1)], X_1, \dots, X_{l-1}, [j, j+(k-1)] = X_l$ form any path of length l in $G_{n,k,b}$, then, for $t = 1, \dots, l$, $\overline{X_t} - \underline{X_{t-1}} \leq b$ and $\overline{X_t} - \underline{X_t} \geq k-1$, which implies

$$\underline{X_t} - \underline{X_{t-1}} \leq b - (k-1). \quad (1)$$

Summing up the inequalities (1) for $t = 1, \dots, l$ yields $j-i \leq l(b-k+1)$. Since $j-i > q(b-k+1)$ we have $l \geq q+1$ and thus the distance is at least $q+1$. \square

Corollary 1. *Let $X, Y \in V_{n,k,b}$ and let $\underline{X} < \underline{Y}$ or $\underline{X} = \underline{Y}$ as well as $\overline{X} < \overline{Y}$. Then X and Y have distance at most $\lceil \frac{\overline{Y} - \underline{X} - b}{b-k+1} \rceil + 1$.*

Proof. Let $i = \underline{X}$ and $j = \overline{Y}$.

Case 1. $j-i \leq b$. Then $\overline{Y} - \underline{X} \leq b$ and $\overline{X} - \underline{Y} \leq \overline{X} - \underline{X} \leq b$. Consequently, X and Y are adjacent by Lemma 1 and their distance is 1. Indeed, from the conditions on X and Y it follows that $\overline{Y} - \underline{X} > k-1$ and thus $\overline{Y} - \underline{X} - b > (-1)(b-k+1)$, which implies $\lceil \frac{\overline{Y} - \underline{X} - b}{b-k+1} \rceil + 1 = 1$.

Case 2. $j-i > b$. Obviously, X is adjacent to $X_1 = [i+b-k+1, i+b]$ and Y is adjacent to $Y_1 = [j-b, j-b+k-1]$.

Case 2.1 $i+b-k+1 \geq j-b$. Then X_1 and Y_1 are adjacent since $\overline{X_1} - \underline{Y_1} = i+b-\underline{Y} = \underline{X} - \underline{Y} + b \leq b$ and $\overline{Y_1} - \underline{X_1} = j-i-b+k-1 \leq b$. Thus X and Y have distance at most 2. Indeed, $2 \leq \lceil \frac{\overline{Y} - \underline{X} - b}{b-k+1} \rceil + 1 \leq \lceil \frac{b+(b-k+1)-b}{b-k+1} \rceil + 1 = 2$.

Case 2.2 $i+b-k+1 < j-b$. By Lemma 2, X_1 and Y_1 have distance $\lceil \frac{j-b-i-(b-k+1)}{b-k+1} \rceil = \lceil \frac{\overline{Y} - \underline{X} - b}{b-k+1} \rceil - 1$, and thus X and Y have distance at most $\lceil \frac{\overline{Y} - \underline{X} - b}{b-k+1} \rceil + 1$. \square

Lemma 3. *The graph $G_{n,k,b}$ has the following number of vertices:*

$$|V_{n,k,b}| = (n-b+1) \binom{b}{k-1} + \binom{b}{k} = (n+1) \binom{b}{k-1} - (k-1) \binom{b+1}{k}.$$

Proof. The proof follows directly from the partition

$$V_{n,k,b} = \bigcup_{i=0}^{n-b} \left\{ X \in \binom{[i, i+b]}{k} : \underline{X} = i \right\} \cup \left\{ X \in \binom{[n-b+1, n]}{k} \right\}.$$

□

4 Bandwidth for $b \geq \frac{n+k-1}{2}$

In the following, we often write the elements of $V_{n,k,b}$ as k -tuples in ascending order, i.e., $X = (i_1, \dots, i_k)$ with $i_1 < \dots < i_k$. Then $\underline{X} = i_1$ and $\overline{X} = i_k$. Furthermore, let $\overleftarrow{X} = (i_k, \dots, i_1)$ as well as $X^c = (n - i_k, \dots, n - i_1)$. We collect all vertices that are adjacent to all other vertices in the set

$$C = \{X \in V_{n,k,b} : n - b \leq \underline{X} \leq \overline{X} \leq b\}.$$

Note that $C \neq \emptyset$ iff $b \geq \frac{n+k-1}{2}$ and that

$$|C| = \binom{2b - n + 1}{k}. \quad (2)$$

We denote the set of remaining vertices by $R = V_{n,k,b} \setminus C$, and split it into two parts:

$$\begin{aligned} R' &= \{X \in R : \underline{X} + \overline{X} \neq n\}, \\ R'' &= \{X \in R : \underline{X} + \overline{X} = n\}. \end{aligned}$$

Let $R'' = R_0'' \cup R_1''$ be a partition of R'' such that

$$\|R_0''\| - \|R_1''\| \leq 1. \quad (3)$$

We define a partition $R' = R_0' \cup R_1'$ of R' by

$$\begin{aligned} R_0' &= \{X \in R' : \underline{X} + \overline{X} < n\}, \\ R_1' &= \{X \in R' : \underline{X} + \overline{X} > n\} \end{aligned}$$

and with $R_0 = R_0' \cup R_0''$ and $R_1 = R_1' \cup R_1''$ we have a partition $R = R_0 \cup R_1$.

Lemma 4. *We have $\|R_0\| - \|R_1\| \leq 1$.*

Proof. A bijection between R_0' and R_1' is given by $X \mapsto X^c$. Hence we have $\|R_0'\| = \|R_1'\|$ and with (3) we obtain the assertion. □

Recall the definition of the *lexicographic ordering* $<_{lex}$ on the set of all k -tuples of integers:

$$(x_1, x_2, \dots, x_k) <_{lex} (y_1, y_2, \dots, y_k) \text{ if } \exists i \in [k] (\forall j \in [i-1] : x_j = y_j) \wedge x_i < y_i. \quad (4)$$

We define a proper numbering of $V_{n,k,b}$ in the form of a total order \leq . The minimal element gets label 1, the next elements get labels 2, 3, ... and the maximal element gets label $|V_{n,k,b}|$. Each total order will be given in the form of an ordinal sum of suborders: If $V_{n,k,b} = S_1 \sqcup \dots \sqcup S_l$ and \leq_i is a total order on S_i , $i = 1, \dots, l$, then $V_{n,k,b} = S_1 \oplus \dots \oplus S_l$ means that the elements of $V_{n,k,b}$ are totally ordered as follows: $X \leq Y$ if there is some i with $X, Y \in S_i$ and $X \leq_i Y$ or there are some i, j with $i < j$ and $X \in S_i$ and $Y \in S_j$. We have $V_{n,k,b} = R_0 \sqcup C \sqcup R_1$, i.e., $l = 3$ with $S_1 = R_0$, $S_2 = C$ and $S_3 = R_1$. We define a total order \leq_{spo} , which we call the *simple palindrom ordering* (SPO), as follows:

$$V_{n,k,b} = R_0 \oplus C \oplus R_1,$$

with the following suborders:

1. For all $X, Y \in R_0$: $X \leq_{spo} Y$ if $X \leq_{lex} Y$.
2. For all $X, Y \in C$: $X \leq_{spo} Y$ if $X \leq_{lex} Y$.
3. For all $X, Y \in R_1$: $X \leq_{spo} Y$ if $\overleftarrow{X} \leq_{lex} \overleftarrow{Y}$.

Let $f_{spo}(X)$ be the label of $X \in V_{n,k,b}$ in the SPO. Recall that the f_{spo} -distance of $X, Y \in V_{n,k,b}$ is given by

$$d_{f_{spo}}(X, Y) = |f_{spo}(X) - f_{spo}(Y)|.$$

Lemma 5. *Let X, Y be two adjacent elements of $V_{n,k,b}$ with $X <_{spo} Y$ and maximal f_{spo} -distance, where in addition \underline{X} is minimal or \overline{Y} is maximal. Then $X = [0, k-1]$ and $Y = [b-k+1, b]$ or $X = [n-b, n-b+k-1]$ and $Y = [n-k+1, n]$.*

Proof. If $Y \in C$ then $X = [0, k-1]$ and $Y = [b-k+1, b]$ have maximal f_{spo} -distance. Analogously, if $X \in C$ then $X = [n-b, n-b+k-1]$ and $Y = [n-k+1, n]$ have maximal f_{spo} -distance. It is not possible that X and Y lie both in R_0 or both in R_1 , because in these cases Y could be replaced by $[b-k+1, b]$ and X by $[n-b, n-b+k-1]$, respectively. Thus it remains the case that $X \in R_0$ and $Y \in R_1$. To reach a maximal f_{spo} -distance, the form $X = (\underline{X}, \underline{X}+1, \dots, \underline{X}+k-1)$ and $Y = (\overline{Y}-k+1, \overline{Y}-k+2, \dots, \overline{Y})$ with $\overline{Y} - \underline{X} = b$ is necessary. So we have $X = [i, i+k-1]$ and $Y = [i+b-k+1, i+b]$ for an $i \in [0, n-b]$. To prove the assertion, it is sufficient to show the following:

- a) For $1 \leq i \leq \frac{n-b}{2}$ we have $d_{spo}([i-1, i+k-2], [i+b-k, i+b-1]) \geq d_{spo}([i, i+k-1], [i+b-k+1, i+b])$.
- b) For $\frac{n-b}{2} < i \leq n-b$ we have $d_{spo}([i-1, i+k-2], [i+b-k, i+b-1]) \leq d_{spo}([i, i+k-1], [i+b-k+1, i+b])$.

We note that b) follows from a) because of the symmetry of the ordering. To show a) we define $I_i = \{X \in V_{n,k,b} : [i, i+k-1] \leq_{spo} X \leq_{spo} [i+b-k+1, i+b]\}$. To prove the inequality it is enough to show that the mapping $X = (i_1, i_2, \dots, i_k) \mapsto (i_1-1, i_2-1, \dots, i_k-1) = \tilde{X}$ is an injection ϕ from I_i to I_{i-1} . The injectivity is clear. Thus it remains to show that $\tilde{X} = \phi(X) \in I_{i-1}$ if $X \in I_i$.

Case 1. $\tilde{X} \in C$.

This case is easy, because $C \subseteq I_j$ for all j .

Case 2. $\tilde{X} \in R_0$.

1. If $X \in R_0$ then $[i, i+k-1] \leq_{lex} X$ and thus $[i-1, i+k-2] \leq_{lex} \tilde{X}$, which yields $\tilde{X} \in I_{i-1}$.
2. If $X \in C$ then $i \leq \frac{n-b}{2} < n-b \leq \underline{X}$ and thus $i-1 < \underline{X}-1 = \underline{\tilde{X}}$. This implies that $[i-1, i+k-2] <_{lex} \tilde{X}$, which yields $\tilde{X} \in I_{i-1}$.
3. If $X \in R_1$ then $\overline{X} \leq \overline{[i+b-k+1, i+b]} = i+b$ because of $X \leq_{spo} [i+b-k+1, i+b]$ and due to $\underline{X} + \overline{X} \geq n$ we have $\underline{X} \geq n-i-b$. This implies $\underline{\tilde{X}} = \underline{X}-1 \geq n-i-b-1 \geq i-1$ because of $i \leq \frac{n-b}{2}$. Hence $\tilde{X} \in I_{i-1}$.

Case 3. $\tilde{X} \in R_1$.

1. The case $X \in R_0$ is not possible because $\underline{\tilde{X}} + \overline{\tilde{X}} = \underline{X} + \overline{X} - 2 < n$, which contradicts $\tilde{X} \in R_1$.
2. Let $X \in C$. Then $\overline{X} \leq b$, which implies $\overline{\tilde{X}} \leq b-1 \leq i+b-2 < \overline{[i+b-k, i+b-1]}$ due to $i \geq 1$. Hence $\overleftarrow{\tilde{X}} <_{lex} \overleftarrow{[i+b-k, i+b-1]}$, which yields $\tilde{X} \in I_{i-1}$.
3. If $X \in R_1$ then $\overline{X} \leq i+b$ and thus $\overline{\tilde{X}} \leq i+b-1$. This implies $\overleftarrow{\tilde{X}} \leq_{lex} \overleftarrow{[i+b-k, i+b-1]}$ and hence $\tilde{X} \in I_{i-1}$.

□

Now the bandwidth of f_{spo} can be determined:

Lemma 6. We have $B_{f_{spo}}(G_{n,k,b}) = \left\lceil \frac{|V_{n,k,b}| + |C| - 2}{2} \right\rceil$.

Proof. We have $f_{spo}([0, k-1]) = 1$ and $f_{spo}([n-k+1, n]) = |V_{n,k,b}|$. If $f_{spo}([n-b, n-b+k-1]) = u+1$, then $f_{spo}([b-k+1, b]) = u+|C|$. Because of Lemma 4 we have $|u - (|V_{n,k,b}| - |C| - u)| \leq 1$ and thus $|(u+|C|-1) - (|V_{n,k,b}| - u - 1)| \leq 1$. Lemma 5 implies that one of the f_{spo} -distances $d_{f_{spo}}([0, k-1], [b-k+1, b]) = u+|C|-1$ and $d_{f_{spo}}([n-b, n-b+k-1], [n-k+1, n]) = |V_{n,k,b}| - u - 1$ is the maximal f_{spo} -distance. As they both differ from each other by at most 1 it follows that

$$B_{f_{spo}}(G_{n,k,b}) = \left\lceil \frac{u+|C|-1 + |V_{n,k,b}| - u - 1}{2} \right\rceil = \left\lceil \frac{|V_{n,k,b}| + |C| - 2}{2} \right\rceil.$$

□

Now we are able to prove the first part of Theorem 1.

Proof of Theorem 1.a). We know from Lemma 6 that $B(G_{n,k,b}) \leq \left\lceil \frac{|V_{n,k,b}| + |C| - 2}{2} \right\rceil$. Let f be an arbitrary proper numbering of $G_{n,k,b}$. Let X_V be the vertex with number 1 and X^V the vertex with number $|V_{n,k,b}|$. Further let X_C be the vertex of C with smallest number, denoted α , and X^C be the vertex of C with largest

number, denoted β . Then $\beta - \alpha \geq |C| - 1$. Further X_V and X^C as well as X_C and X^V are adjacent with $d_f(X_V, X^C) = \beta - 1$ and $d_f(X_C, X^V) = |V_{n,k,b}| - \alpha$. The sum of them is

$$s = (\beta - 1) + (|V_{n,k,b}| - \alpha) = |V_{n,k,b}| + (\beta - \alpha) - 1 \geq |V_{n,k,b}| + |C| - 2.$$

The maximum of both f -distances is therefore at least $\left\lceil \frac{|V_{n,k,b}| + |C| - 2}{2} \right\rceil$. From Lemma 3 and (2) we obtain

$$B(G_{n,k,b}) = \left\lceil \frac{(n+1)\binom{b}{k-1} - (k-1)\binom{b+1}{k} + \binom{2b-n+1}{k} - 2}{2} \right\rceil.$$

□

5 Asymptotic bandwidth for $b = o(n)$

In this section, we consider the case, where b grows sublinearly with respect to n . First we take a simple proper numbering, which provides an upper bound for the bandwidth.

Lemma 7. *Let n, k, b be arbitrary integers with $1 \leq k-1 \leq b \leq n$. Then*

$$B(G_{n,k,b}) \leq k \binom{b}{k}.$$

Proof. We order the vertices of $G_{n,k,b}$ in a lexicographic way, see (4). Let $f_{lex}(X)$ be the label of $X \in V_{n,k,b}$ with respect to this ordering. Now let X and Y be two adjacent vertices with $X <_{lex} Y$ and let $X' = [\underline{X}, \underline{X} + k - 1]$ and $Y' = [\underline{X} + b - k + 1, \underline{X} + b]$. Then

$$X' \leq_{lex} X <_{lex} Y \leq_{lex} Y'. \quad (5)$$

Moreover, for $j \in [0, n - b]$,

$$|\{X \in V_{n,k,b} : \underline{X} = j\}| = \binom{b}{k-1}.$$

and, for $j \in [n - b + 1, n]$,

$$|\{X \in V_{n,k,b} : \underline{X} = j\}| \leq \binom{b}{k-1}.$$

Since Y' is the lexicographically smallest vertex with minimum element $\underline{X} + b - k + 1$ it follows that

$$|f_{lex}(Y') - f_{lex}(X')| \leq (b - k + 1) \binom{b}{k-1} = k \binom{b}{k}. \quad (6)$$

Now (5) and (6) imply

$$|f_{lex}(Y) - f_{lex}(X)| \leq k \binom{b}{k},$$

which proves the assertion. \square

Chvátal observed in [6] that a lower bound for the bandwidth is given by

$$B(G) \geq \left\lceil \frac{|V| - 1}{\text{diam}(G)} \right\rceil. \quad (7)$$

Here the *diameter* $\text{diam}(G)$ of the graph $G = (V, E)$ is the maximal distance of any two vertices of G .

By Corollary 1, the distance of any two vertices of $G_{n,k,b}$ is at most $\lceil \frac{n-0-b}{b-k+1} \rceil + 1 = \lceil \frac{n-k+1}{b-k+1} \rceil$ and by Lemma 2 the vertices $[0, k-1]$ and $[n-k+1, n]$ have distance $\lceil \frac{n-k+1}{b-k+1} \rceil$. Accordingly,

$$\text{diam}(G_{n,k,b}) = \left\lceil \frac{n-k+1}{b-k+1} \right\rceil. \quad (8)$$

Now we have all preparations to prove Theorem 1 b).

Proof of Theorem 1 b). From Lemma 7 we know that

$$B(G_{n,k,b}) \leq k \binom{b}{k}.$$

For the lower bound, we use the fact that $\binom{b}{k} = \frac{b^k}{k!} + O(b^{k-1})$ as $b \rightarrow \infty$. We have by Lemma 3, (7) and (8) for $n \rightarrow \infty$:

$$\begin{aligned} B(G_{n,k,b}) &\geq \left\lceil \frac{|V_{n,k,b}| - 1}{\text{diam}(G_{n,k,b})} \right\rceil = \frac{(n-b+1)\binom{b}{k-1} + \binom{b}{k} - 1}{\lceil \frac{n-k+1}{b-k+1} \rceil} = \frac{n\binom{b}{k-1} + O(b^k)}{\frac{n-k+1}{b-k+1} + O(1)} \\ &= \frac{nk\binom{b}{k} + O(b^{k+1})}{n + O(b)} = \frac{k\binom{b}{k} + O(\frac{b^{k+1}}{n})}{1 + O(\frac{b}{n})}. \end{aligned}$$

In the case $b = o(n^{\frac{1}{k+1}})$ we have

$$\frac{k\binom{b}{k} + O(\frac{b^{k+1}}{n})}{1 + O(\frac{b}{n})} \geq \left(k\binom{b}{k} + o(1) \right) \left(1 - O(\frac{b}{n}) \right) = k\binom{b}{k} - o(1).$$

Thus, $B(G_{n,k,b}) \geq k\binom{b}{k}$ for sufficiently large n , which proves the first part of the assertion.

If $b = o(n)$ then

$$\frac{k\binom{b}{k} + O(\frac{b^{k+1}}{n})}{1 + O(\frac{b}{n})} \sim k\binom{b}{k},$$

which shows that $B(G_{n,k,b}) \gtrsim k\binom{b}{k}$ as $n \rightarrow \infty$, which proves the second part of the assertion. \square

6 Further basic properties for the asymptotics

Lemma 8. *Let $b \sim \beta n$, $\delta > 0$, $X, Y \in V_{n,k,b}$ and $\overline{Y} - \underline{X} \lesssim i(1-\delta)\beta n$ as well as $\overline{X} - \underline{Y} \lesssim i(1-\delta)\beta n$, where $n \rightarrow \infty$ and i is a positive integer. If n is sufficiently large, then X and Y have distance at most i .*

Proof. Without loss of generality, let $(\underline{X}, \overline{X}) \leq_{lex} (\underline{Y}, \overline{Y})$, i.e., $\underline{X} < \underline{Y}$ or $\underline{X} = \underline{Y}$ as well as $\overline{X} < \overline{Y}$. By Corollary 1, X and Y have distance at most $\lceil \frac{i(1-\delta)\beta n - b}{b-k+1} \rceil + 1 = \lceil \frac{i(1-\delta)\beta - \beta + o(1)}{\beta + o(1)} \rceil + 1 = \lceil i(1-\delta) - 1 + o(1) \rceil + 1 \leq i$. \square

Let P be a polygon in $\Omega = \{(x, y) \in \mathbb{R}^2 : 0 \leq x \leq y \leq 1\}$ and let $\text{int}(P)$ be the interior of P . Let

$$V_{n,k}(P) = \left\{ X \in \binom{[0, n]}{k} : \frac{1}{n}(\underline{X}, \overline{X}) \in P \right\},$$

$$V_{n,k}^o(P) = \left\{ X \in \binom{[0, n]}{k} : \frac{1}{n}(\underline{X}, \overline{X}) \in \text{int}(P) \right\}.$$

Lemma 9. *We have*

$$|V_{n,k}^o(P)| \sim |V_{n,k}(P)| \sim \left(\frac{1}{(k-2)!} \iint_P (y-x)^{k-2} dx dy \right) n^k \text{ as } n \rightarrow \infty.$$

Proof. Let i, j be integers with $0 \leq i \leq j \leq n$. Obviously,

$$|\{X \in \binom{[0, n]}{k} : \underline{X} = i, \overline{X} = j\}| = \binom{j-i-1}{k-2}.$$

Thus

$$|V_{n,k}(P)| = \sum_{\frac{1}{n}(i,j) \in P} \binom{j-i-1}{k-2} = \sum_{\frac{1}{n}(i,j) \in P} \binom{n(j/n - i/n) - 1}{k-2}.$$

For $0 \leq z \leq 1$, we have

$$\binom{nz-1}{k-2} \leq \frac{z^{k-2}}{(k-2)!} n^{k-2}.$$

Accordingly,

$$\frac{1}{n^k} |V_{n,k}(P)| \leq \frac{1}{(k-2)!} \sum_{\frac{1}{n}(i,j) \in P} \frac{1}{n^2} (j/n - i/n)^{k-2}.$$

The RHS is a Riemann sum for the integral $\iint_P (y-x)^{k-2} dx dy$, which shows that

$$|V_{n,k}(P)| \lesssim \left(\frac{1}{(k-2)!} \iint_P (y-x)^{k-2} dx dy \right) n^k \text{ as } n \rightarrow \infty.$$

For $\delta > 0$ let $P_\delta = \{(x, y) \in P : y-x \geq \delta\}$. Clearly, $P_\delta \subseteq P$. Obviously, for any $\varepsilon > 0$ there is some $\delta > 0$ such that

$$\iint_{P_\delta} (y-x)^{k-2} dx dy \geq (1-\varepsilon) \iint_P (y-x)^{k-2} dx dy.$$

Moreover, for any $z \geq \delta$ there is some n_0 such that for all $n > n_0$

$$\frac{k-2}{nz} \leq \varepsilon.$$

This implies

$$(1-\varepsilon)^{k-2} \frac{(nz)^{k-2}}{(k-2)!} \leq \left(1 - \frac{k-2}{nz}\right)^{k-2} \frac{(nz)^{k-2}}{(k-2)!} \leq \binom{nz-1}{k-2}$$

and further

$$\begin{aligned} \frac{1}{n^k} |V_{n,k}(P)| &\geq (1-\varepsilon)^{k-2} \frac{1}{(k-2)!} \sum_{\frac{1}{n}(i,j) \in P_\delta} \frac{1}{n^2} (j/n - i/n)^{k-2} \\ &\gtrsim (1-\varepsilon)^{k-1} \frac{1}{(k-2)!} \iint_P (y-x)^{k-2} dx dy. \end{aligned}$$

Now, with $\varepsilon \rightarrow 0$ we obtain

$$|V_{n,k}(P)| \gtrsim \left(\frac{1}{(k-2)!} \iint_P (y-x)^{k-2} dx dy \right) n^k \text{ as } n \rightarrow \infty.$$

The reasoning for $|V_{n,k}^o(P)|$ is the same. □

For the sake of brevity, we define the *measure of the polygon* $P \subseteq \Omega$ by

$$\mu(P) = \frac{1}{(k-2)!} \iint_P (y-x)^{k-2} dx dy.$$

Corollary 2. *If $S \subseteq \binom{[0,n]}{k}$ is a family of sets that contains all $X \in \binom{[0,n]}{k}$ with $\frac{1}{n}(\underline{X}, \overline{X}) \in \text{int}(P)$ and some $X \in \binom{[0,n]}{k}$ with $\frac{1}{n}(\underline{X}, \overline{X})$ on the boundary of P , then*

$$|S| \sim \mu(P) n^k \text{ as } n \rightarrow \infty.$$

7 Definition and measure of crucial polygons

Recall that we consider the case $b \sim \beta n$ with $\beta \in (0, 0.5]$ and $q \in \mathbb{N}$ and $r \in \mathbb{R}$ such that $1 = q\beta + r$. We define in Ω several sets of points. First let for $i = 1, \dots, q$

$$\begin{aligned} A_i &= (ir, ir + q(\beta - r)), \\ B_i &= (r - \beta + i\beta, r - \beta + i\beta), \\ C_i &= (i\beta, i\beta). \end{aligned}$$

In the following we denote lines given by $y = ax + b$ (or, more generally, by $ax + by = c$) by $g_{y=ax+b}$ (or $g_{ax+by=c}$). Note that the points B_i and C_i lie on $g_{y=x}$. Furthermore, the points A_i lie above or on the line $g_{y=x+\beta}$ iff $r \leq \frac{q-1}{q^2+q-1}$. This is the reason for the distinction between a) and b) in Theorem 2.

For the Case a), i.e., for $r \leq \frac{q-1}{q^2+q-1}$, we define points D_i and E_i as the intersection points of the segments $\overline{A_i B_i}$ resp. $\overline{A_i C_i}$ with the line $g_{y=x+\beta}$. An easy computation yields with

$$\gamma = \beta(1 - 1/q)$$

that for $i = 1, \dots, q$

$$\begin{aligned} D_i &= (r + (i-1)\gamma, r + (i-1)\gamma + \beta), \\ E_i &= (i\gamma, i\gamma + \beta). \end{aligned}$$

For the Case b), i.e., for $r > \frac{q-1}{q^2+q-1}$, we define points F_i as those points that provide an equipartition into $q+1$ parts of the segment between $(0,0)$ and $(1,1)$ and points G_i as the intersection points of the segments between F_i and $(0,1)$ with the line $g_{y=x+\beta}$. It is straightforward that for $i = 0, \dots, q+1$

$$\begin{aligned} F_i &= (i/(q+1), i/(q+1)), \\ G_i &= (i(1-\beta)/(q+1), i(1-\beta)/(q+1) + \beta). \end{aligned}$$

As for the Case a), let E_q be the intersection point of the lines $g_{y=x+\beta}$ and $g_{y=q\beta}$. Finally, we define an auxiliary point $H_1 = (r, \beta)$. For the Case a) we also use the points $C_0 = F_0 = (0,0)$, $B_{q+1} = F_{q+1} = (1,1)$, $E_0 = G_0 = (0, \beta)$ and $D_{q+1} = G_{q+1} = (1-\beta, 1)$.

In Figures 1 and 2 the points are illustrated for $\beta = 9/20$ (i.e., Case a) with $q = 2, r = 1/10 < 1/5$) as well as for $\beta = 7/20$ (i.e., Case b) with $q = 2, r = 3/10 > 1/5$).

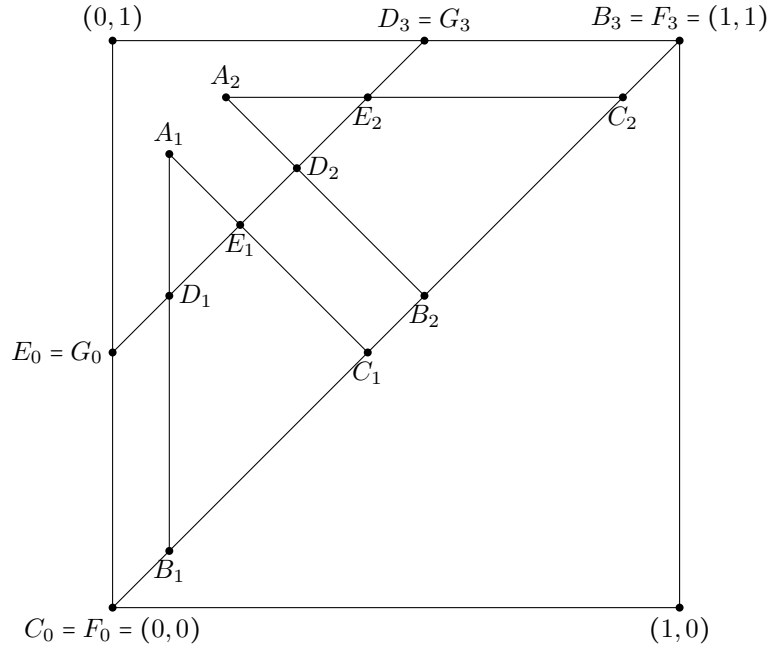


Figure 1: Polygons and important points for $\beta = 9/20$ (i.e., Case a) with $q = 2, r = 1/10 < 1/5$).

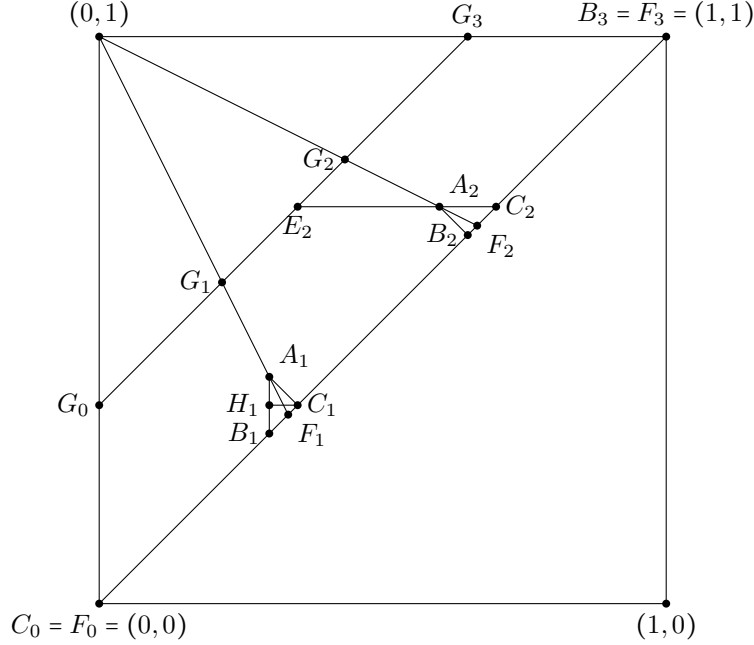


Figure 2: Polygons and important points for $\beta = 7/20$ (i.e., Case b) with $q = 2, r = 3/10 > 1/5$).

It is easy to check that the points A_i lie on the segments $\overrightarrow{F_i G_i}$ (we have $\overrightarrow{F_i A_i} = (1 - r(q+1)/\beta)\overrightarrow{F_i G_i}$) and that the points $F_0, B_1, F_1, C_1, \dots, B_q, F_q, C_q, F_{q+1}$ lie in this order on the line $g_{y=x}$. It is also easy to see that H_1 lies on the segment $\overrightarrow{A_1 B_1}$.

Lemma 10. *Let $0 \leq s < t \leq 1$, let $P_1 = (\xi_1, \xi_1 + s), P_2 = (\xi_2, \xi_2 + s)$ be points on the line $g_{y=x+s}$ with $0 \leq \xi_1 \leq \xi_2 \leq 1 - s$ and let $P_3 = (\xi_3, \xi_3 + t), P_4 = (\xi_4, \xi_4 + t)$ be points on the line $g_{y=x+t}$ with $0 \leq \xi_3 \leq \xi_4 \leq 1 - t$. Let $u = \xi_2 - \xi_1$ and $v = \xi_4 - \xi_3$. Then*

$$\mu(P_1 P_2 P_4 P_3) = \frac{1}{(k-2)!} \frac{1}{t-s} \left(\frac{1}{k} (v-u)(t^k - s^k) + \frac{tu - sv}{k-1} (t^{k-1} - s^{k-1}) \right).$$

Proof. The proof follows directly by computing the integrals, using the coordinate transformation $x' = x + y, y' = y - x$. This leads to a domain of integration in form of a trapezoid whose basis is parallel to the x' -axis. \square

Since only the difference of the ξ -values has influence we have:

Corollary 3. *Let $0 \leq s < t \leq 1$, let P_1, P_2 and Q_1, Q_2 be points of Ω on the line $g_{y=x+s}$ with $\overrightarrow{P_1 P_2} = \overrightarrow{Q_1 Q_2}$ showing to north east and let P_3, P_4 and Q_3, Q_4 be points on the line $g_{y=x+t}$ with $\overrightarrow{P_3 P_4} = \overrightarrow{Q_3 Q_4}$ showing to north east. Then*

$$\mu(P_1 P_2 P_4 P_3) = \mu(Q_1 Q_2 Q_4 Q_3).$$

By inserting the corresponding values and using the definition of the c -functions in Theorem 2, we obtain:

Corollary 4. *We have for all possible i :*

$$\begin{aligned}
\mu(F_0 F_{q+1} G_{q+1} G_0) &= \frac{\beta^{k-1}}{k!} (k - (k-1)\beta) = (q+1)c_2(\beta, k), \\
\mu(C_i B_{i+1} D_{i+1} E_i) &= \frac{\beta^{k-1}}{(k-1)!} r = (q+1)c_2(\beta, k) - qc_1(\beta, k), \\
\mu(B_i C_i E_i D_i) &= \frac{\beta^{k-1}}{k!} \left(k(\beta - r) - \frac{\beta(k-1)}{q} \right) \\
&= (q+1)(c_1(\beta, k) - c_2(\beta, k)), \\
\mu(F_i F_{i+1} G_{i+1} G_i) &= \frac{1}{q+1} \mu(F_0 F_{q+1} G_{q+1} G_0) = c_2(\beta, k), \\
\mu(A_i B_i C_i) &= \frac{(\beta - r)^k}{k!} q^{k-1} = (q+1)c_3(\beta, k), \\
\mu(A_i B_i F_i) &= \left(1 - \frac{i}{q+1}\right) \mu(A_i B_i C_i) = (q+1-i)c_3(\beta, k), \\
\mu(A_i F_i C_i) &= \frac{i}{q+1} \mu(A_i B_i C_i) = ic_3(\beta, k), \\
\mu(B_1 C_1 H_1) &= \frac{1}{q^{k-1}} \mu(A_1 B_1 C_1) = \frac{q+1}{q^{k-1}} c_3(\beta, k).
\end{aligned}$$

8 Proof of the lower bounds for the bandwidth in Theorem 2

We fix some $\varepsilon > 0$, where ε is sufficiently small, in particular $\varepsilon < 1$. Let P be a polygon in Ω . In the following we work with dilations of P around a given centerpoint by factors of the form $(1 - \varepsilon)$ and $(1 + \varepsilon)$, respectively. We denote the new polygons by \underline{P} and \overline{P} , respectively. But we emphasize that \underline{P} and \overline{P} depend on ε and on the centerpoint of dilation. Note that

$$\lim_{\varepsilon \rightarrow 0} \mu(\underline{P}) = \lim_{\varepsilon \rightarrow 0} \mu(\overline{P}) = \mu(P).$$

Though the following result follows also directly from Lemma 3, we prove it as an example for our polygon-method:

Lemma 11. *We have*

$$|V_{n,k,b}| \sim \mu(F_0 F_{q+1} G_{q+1} G_0) n^k = (q+1)c_2(\beta, k) n^k.$$

Proof. Let briefly $T = F_0 F_{q+1} G_{q+1} G_0$. Note that $T = \{(x, y) \in \Omega : y \leq x + \beta\}$. We choose F_0 as the centerpoint of dilation so that e.g. the vertices of \underline{T} are given by $\underline{F}_0 = (0, 0)$, $\underline{F}_{q+1} = (1 - \varepsilon)(1, 1)$, $\underline{G}_{q+1} = (1 - \varepsilon)(1 - \beta, 1)$, $\underline{G}_0 = (1 - \varepsilon)(0, \beta)$. If the dilation factors are $(1 - \varepsilon)$ and $(1 + \varepsilon)$, respectively, then $\underline{T} \subseteq \{(x, y) \in \Omega : y \leq$

$x + (1 - \varepsilon)\beta\}$ and $\overline{T} \cap \Omega \subseteq \{(x, y) \in \Omega : y \leq x + (1 + \varepsilon)\beta\}$. It is easy to check that for sufficiently large n and any $X \in \binom{[0, n]}{k}$ the following implications are true:

$$\frac{1}{n}(\underline{X}, \overline{X}) \in \underline{T} \Rightarrow X \in V_{n,k,b} \Rightarrow \frac{1}{n}(\underline{X}, \overline{X}) \in \overline{T} \cap \Omega.$$

Thus

$$|V_{n,k}(\underline{T})| \leq |V_{n,k,b}| \leq |V_{n,k}(\overline{T} \cap \Omega)|$$

and hence by Corollary 2

$$\mu(\underline{T})n^k \lesssim |V_{n,k,b}| \lesssim \mu(\overline{T})n^k.$$

With $\varepsilon \rightarrow 0$ we obtain (using also Corollary 4)

$$|V_{n,k,b}| \sim \mu(T)n^k = (q+1)c_2(\beta, k)n^k.$$

□

Now we prove the first asymptotic lower bound:

Lemma 12. *We have*

$$B(G_{n,k,b}) \gtrsim \frac{1}{q}\mu(F_0 C_q E_q G_0)n^k = c_1(\beta, k)n^k \text{ as } n \rightarrow \infty.$$

Proof. Let briefly $Q = F_0 C_q E_q G_0$. For the dilation, we choose again F_0 as the centerpoint. With the factor $(1 - \varepsilon)$ we obtain \underline{Q} . Let G' be the subgraph of $G_{n,k,b}$ induced by

$$V' = \left\{ X \in V_{n,k,b} : \frac{1}{n}(\underline{X}, \overline{X}) \in \underline{Q} \right\}.$$

Note that $\frac{1}{n}\overline{X} \leq q\beta(1 - \varepsilon)$ for all $X \in V'$.

Now we show that for sufficiently large n

$$\text{diam}(G') \leq q. \tag{9}$$

Let X and Y be any two distinct vertices of G' . Then $0 \leq \underline{X} \leq \overline{X} \leq q\beta(1 - \varepsilon)n$ and $0 \leq \underline{Y} \leq \overline{Y} \leq q\beta(1 - \varepsilon)n$. By Lemma 8, X and Y have distance at most q if n is sufficiently large, which proves (9). From the Chvátal bound (7) it follows that

$$B(G_{n,k,b}) \geq B(G') \geq \frac{|V'| - 1}{q} \gtrsim \frac{\mu(\underline{Q})}{q}n^k$$

and with $\varepsilon \rightarrow 0$

$$B(G_{n,k,b}) \gtrsim \frac{\mu(Q)}{q}n^k.$$

Let briefly $Q_q = C_q B_{q+1} D_{q+1} E_q = C_q F_{q+1} G_{q+1} E_q$ and recall $T = F_0 F_{q+1} G_{q+1} G_0$. Then $Q = T \setminus Q_q$ and hence

$$\mu(Q) = \mu(T) - \mu(Q_q).$$

By Corollary 4,

$$\mu(Q) = (q+1)c_2(\beta, k) - ((q+1)c_2(\beta, k) - qc_1(\beta, k)) = qc_1(\beta, k).$$

□

Now we prove the second asymptotic lower bound, which applies only for the second case:

Lemma 13. *If $r > \frac{q-1}{q^2+q-1}$ then*

$$\begin{aligned} B(G_{n,k,b}) &\gtrsim \left(\mu(F_0 F_1 G_1 G_0) + \frac{1}{q+1} \mu(B_1 C_1 H_1) \right) n^k \\ &= \left(c_2(\beta, k) + \frac{1}{q^{k-1}} c_3(\beta, k) \right) n^k. \end{aligned}$$

Proof. Let briefly $R = B_1 C_1 H_1$. We choose $\frac{1}{2}(B_1 + C_1) = \frac{1}{2}(r + \beta, r + \beta)$ as the centerpoint and $(1 - \frac{2}{\beta-r}\varepsilon)$ as the factor of dilation and thus obtain \underline{R} from R . The vertices of \underline{R} are $\underline{B}_1 = (r + \varepsilon, r + \varepsilon)$, $\underline{C}_1 = (\beta - \varepsilon, \beta - \varepsilon)$ and $\underline{H}_1 = (r + \varepsilon, \beta - \varepsilon)$. Let f be a bandwidth numbering of $G_{n,k,b}$. Let X_V and X^V be those vertices for which $f(X_V) = 1$ and $f(X^V) = |V_{n,k,b}|$.

Case 1. $\underline{X}_V \leq n(r + \varepsilon)$ and $\underline{X}^V \leq n(r + \varepsilon)$.

Then, for sufficiently small ε and sufficiently large n , $\overline{X}_V, \overline{X}^V \leq n(r + \varepsilon + \beta + \varepsilon) \leq 2\beta(1 - \varepsilon)n$. Lemma 8 implies that X_V and X^V have distance at most 2 and by Lemma 11, we have

$$B(G_{n,k,b}) \geq \frac{|V_{n,k,b}| - 1}{2} \gtrsim \frac{q+1}{2} c_2(\beta, k) n^k \gtrsim \left(c_2(\beta, k) + \frac{1}{q^{k-1}} c_3(\beta, k) \right) n^k.$$

Case 2. $\underline{X}_V > n(r + \varepsilon)$ and $\underline{X}^V > n(r + \varepsilon)$.

Then, for sufficiently large n , $\overline{X}^V - \overline{X}_V, \overline{X}_V - \overline{X}^V \leq n - n(r + \varepsilon) = (1 - r - \varepsilon)n = q\beta(1 - \frac{\varepsilon}{q\beta})n$. Lemma 8 implies that X_V and X^V have distance at most q and by Lemma 11, we have

$$B(G_{n,k,b}) \geq \frac{|V_{n,k,b}| - 1}{q} \gtrsim \frac{q+1}{q} c_2(\beta, k) n^k \gtrsim \left(c_2(\beta, k) + \frac{1}{q^{k-1}} c_3(\beta, k) \right) n^k.$$

Case 3. $\underline{X}_V \leq n(r + \varepsilon)$ and $\underline{X}^V > n(r + \varepsilon)$.

Let $X_R = \arg \min\{f(X) : X \in V_{n,k}(\underline{R})\}$ and $X^R = \arg \max\{f(X) : X \in V_{n,k}(\underline{R})\}$. Clearly,

$$f(X^R) - f(X_R) \geq |V_{n,k}(\underline{R})| - 1. \quad (10)$$

In view of $n(r + \varepsilon) \leq \overline{X}_R \leq n(\beta - \varepsilon)$, $n(r + \varepsilon) \leq \overline{X}_V \leq n(r + \varepsilon + \beta + \varepsilon)$ and $r < \beta$ we have $\overline{X}_R - \underline{X}_V \leq n(\beta - r - 2\varepsilon) \leq b$ and $\overline{X}_V - \underline{X}_R \leq n(\beta + \varepsilon) \leq b$ if ε is sufficiently small and n is sufficiently large. Thus, by Lemma 1, X_V and X_R are adjacent and hence

$$f(X^R) - f(X_V) \leq B(G_{n,k,b}). \quad (11)$$

In view of $\underline{X}_R, \underline{X}^V \geq n(r + \varepsilon)$ and analogously to Case 2, X_R and X^V have distance at most q and hence

$$f(X^V) - f(X_R) \leq qB(G_{n,k,b}). \quad (12)$$

From (10), (11) and (12) we obtain

$$f(X^V) - f(X_V) \leq (q+1)B(G_{n,k,b}) - (|V_{n,k}(\underline{R})| - 1)$$

and thus by Corollary 4 and Lemma 11

$$B(G_{n,k,b}) \gtrsim c_2(\beta, k)n^k + \left(1 - \frac{2}{\beta - r}\varepsilon\right)^2 \frac{1}{q^{k-1}}c_3(\beta, k)n^k$$

and with $\varepsilon \rightarrow 0$ the assertion follows. \square

9 Proof of the upper bounds for the bandwidth in Theorem 2

In the following, we present two proper numberings f of $G_{n,k,b}$ whose bandwidth is asymptotically equal to the asserted upper bounds. As in Section 4, we define a total order $V_{n,k,b} = S_1 \oplus \dots \oplus S_l$ with suborders given by means of polygons. In order to avoid intersections on the boundaries we explicitly describe which part of the boundary is deleted, though the ordering can be given on the whole polygon. For example, a notation of the form $C_0B_1D_1E_0 \setminus \overline{B_1D_1}$ means that the segment $\overline{B_1D_1}$ is deleted from the closed quadrangle $C_0B_1D_1E_0$.

Case a) $r \leq \frac{q-1}{q^2+q+1}$.

We define the total order \leq as follows:

$$\begin{aligned} V_{n,k,b} = & V_{n,k}(C_0B_1D_1E_0 \setminus \overline{B_1D_1}) \oplus V_{n,k}(B_1C_1E_1D_1 \setminus \overline{C_1E_1}) \oplus V_{n,k}(C_1B_2D_2E_1 \setminus \overline{B_2D_2}) \\ & \oplus \dots \oplus V_{n,k}(B_qC_qE_qD_q \setminus \overline{C_qE_q}) \oplus V_{n,k}(C_qB_{q+1}D_{q+1}E_q). \end{aligned}$$

We still have to define the ordering of the elements of $V_{n,k}(C_iB_{i+1}D_{i+1}E_i)$, $i = 0, \dots, q$, and of $V_{n,k}(B_iC_iE_iD_i)$, $i = 1, \dots, q$ (here we may allow the complete boundary).

If $(\underline{X}, \overline{X}) = (\underline{Y}, \overline{Y})$ we set in both cases $X \leq_i Y$ if $X \leq_{lex} Y$. Thus let $(\underline{X}, \overline{X}) \neq (\underline{Y}, \overline{Y})$.

First we discuss $V_{n,k}(C_iB_{i+1}D_{i+1}E_i)$. We use a new coordinate system with the same origin and with transformation matrix and inverse transformation matrix

$$M_i = \begin{pmatrix} 1 & -i/q \\ 1 & 1 - i/q \end{pmatrix} \text{ and } M_i^{-1} = \begin{pmatrix} 1 - i/q & i/q \\ -1 & 1 \end{pmatrix}.$$

Thus the new coordinate axes have direction of $\overrightarrow{C_iB_{i+1}}$ and $\overrightarrow{C_iE_i}$. The ordering is a lexicographic ordering of the points $\frac{1}{n}(\underline{X}, \overline{X})$ with respect to the new

coordinate system, i.e., for $X, Y \in V_{n,k}(C_i B_{i+1} D_{i+1} E_i)$ and $(\underline{X}, \overline{X}) \neq (\underline{Y}, \overline{Y})$ we set

$$X \leq_i Y \text{ if } ((1 - i/q)\underline{X} + i\overline{X}/q, -\underline{X} + \overline{X}) \leq_{lex} ((1 - i/q)\underline{Y} + i\overline{Y}/q, -\underline{Y} + \overline{Y}),$$

see Figure 3.

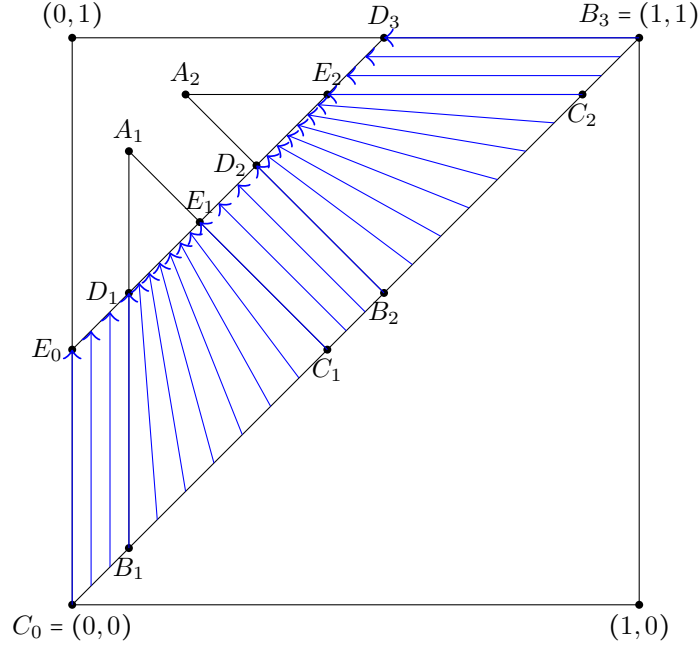


Figure 3: Schematic illustration of the ordering for Case a).

Now we discuss $V_{n,k}(B_i C_i E_i D_i)$. Here we work with polar coordinates in the coordinate system with origin A_i and x -axis in the direction of $\overrightarrow{A_i B_i}$ and arbitrary, but fixed unit length. For $X \in V_{n,k}(B_i C_i E_i D_i)$ let $\varphi_i(X)$ and $r_i(X)$ be the angular and radial coordinates of $\frac{1}{n}(\underline{X}, \overline{X})$ in this coordinate system. The ordering is a lexicographic ordering with respect to the *reflected* polar coordinates, i.e., for $X, Y \in V_{n,k}(B_i C_i E_i D_i)$ and $(\underline{X}, \overline{X}) \neq (\underline{Y}, \overline{Y})$ we set

$$X \leq_i Y \text{ if } (\varphi_i(X), -r_i(X)) \leq_{lex} (\varphi_i(Y), -r_i(Y)),$$

see Figure 3.

Note that for simpler numerical computations $\varphi_i(X)$ may be enlarged to an angle such that one leg is parallel to the y -axis, the size of the angle may be replaced by $\tan(\varphi_i(X))$ and the Euclidean norm for $r_i(X)$ may be replaced by some other norm, e.g. the L_1 -norm.

It is easy to check that

$$X \leq Y \text{ implies } \underline{X} \leq \overline{Y}. \quad (13)$$

Lemma 14. *Let f be the numbering for Case a). Then, for $n \rightarrow \infty$,*

$$B_f(G) \lesssim c_1(\beta, k)n^k.$$

Proof. Let $P = (\xi, \xi)$ be any point on the segment $\overline{C_0 B_{q+1}}$, i.e., $0 \leq \xi \leq 1$. With P we associate a new point \hat{P} as follows: If $P \in \overline{C_i B_{i+1}}$ for some i then let \hat{P} be the intersection point of the line $g_{y=x+\beta}$ with the line through P that is parallel to $C_i E_i$. If $P \in \overline{B_i C_i}$ for some i then let \hat{P} be the intersection point of the line $g_{y=x+\beta}$ with the line through P and A_i , see Figure 4.

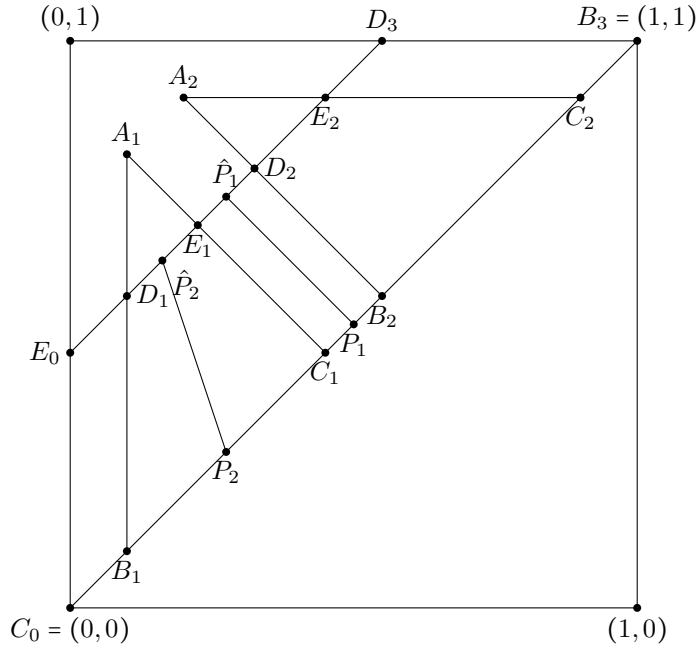


Figure 4: Important points for the upper bound for Case a).

Moreover, for $X \in V_{n,k,b}$ let

$$P_{\underline{X}} = \frac{1}{n}(\underline{X}, \underline{X}) \text{ and } P_{\overline{X}} = \frac{1}{n}(\overline{X}, \overline{X}). \quad (14)$$

Let $X, Y \in V_{n,k,b}$ with $X \leq Y$. By the definition of the ordering and in view of (13),

$$f(Y) - f(X) \leq |V_{n,k}(P_{\underline{X}} P_{\overline{Y}} \hat{P}_{\overline{Y}} \hat{P}_{\underline{X}})|.$$

Thus we have to prove that

$$|V_{n,k}(P_X P_{\bar{Y}} \hat{P}_{\bar{Y}} \hat{P}_X)| \lesssim c_1(\beta, k) n^k. \quad (15)$$

Note that

$$\overrightarrow{B_i B_{i+1}} = \overrightarrow{C_i C_{i+1}} = (\beta, \beta). \quad (16)$$

Since X and Y are adjacent we have by Lemma 1, $\overline{Y} - \underline{X} \leq b \sim \beta n$. Let $\varepsilon > 0$. Then $\frac{1}{n}(\overline{Y} - \underline{X}) \leq \beta + \varepsilon$ for sufficiently large n . Let

$$P' = \left(\frac{1}{n}\underline{X} + \beta, \frac{1}{n}\underline{X} + \beta\right). \quad (17)$$

If $\beta < \frac{1}{n}(\overline{Y} - \underline{X}) \leq \beta + \varepsilon$ then $\overrightarrow{P'P_Y} = (\delta, \delta)$ with $\delta \leq \varepsilon$. In view of

$$\mu(P_{\underline{X}}P_{\overline{Y}}\hat{P}_{\overline{Y}}\hat{P}_{\underline{X}}) = \mu(P_{\underline{X}}P'\hat{P}'\hat{P}_{\underline{X}}) + \mu(P'P_{\overline{Y}}\hat{P}_{\overline{Y}}\hat{P}')$$

we have

$$\mu(P_{\underline{X}}P_{\overline{Y}}\hat{P}_{\overline{Y}}\hat{P}_{\underline{X}}) \leq \mu(P_{\underline{X}}P'\hat{P}'\hat{P}_{\underline{X}}) + O(\varepsilon),$$

which is clearly also true if $\frac{1}{n}(\overline{Y} - \underline{X}) \leq \beta$. Using Corollary 4, it is easy to check that for all possible i ,

$$\begin{aligned} \mu(P_{\underline{X}}P'\hat{P}'\hat{P}_{\underline{X}}) &= \mu(B_i B_{i+1} D_{i+1} D_i) = \mu(C_i C_{i+1} E_{i+1} E_i) \\ &= \mu(B_i C_i E_i D_i) + \mu(C_i B_{i+1} D_{i+1} E_i) \\ &= c_1(\beta, k). \end{aligned} \quad (18)$$

An illustration of this fact can be found in Figures 5 and 6.

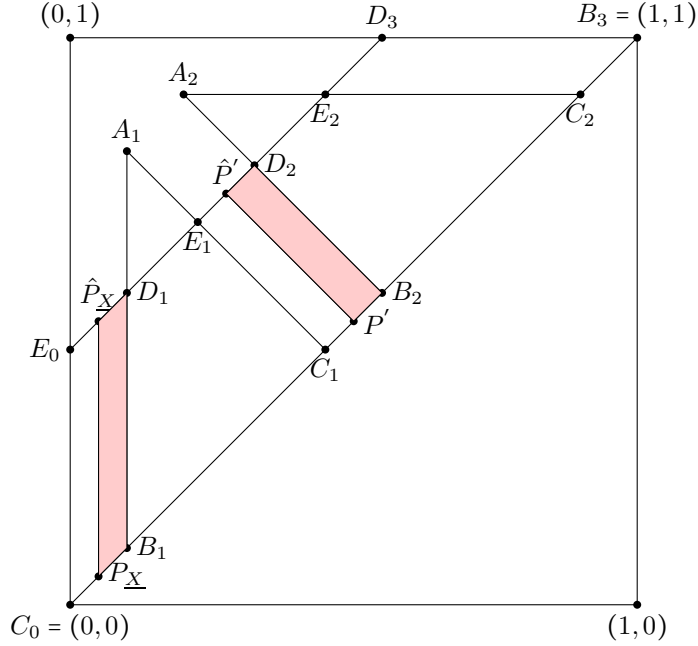


Figure 5: Illustration of (18), where $P_{\underline{X}} \in \overline{C_i B_{i+1}}$ for some i . Both red quadrangles have the same measure by Corollary 3.

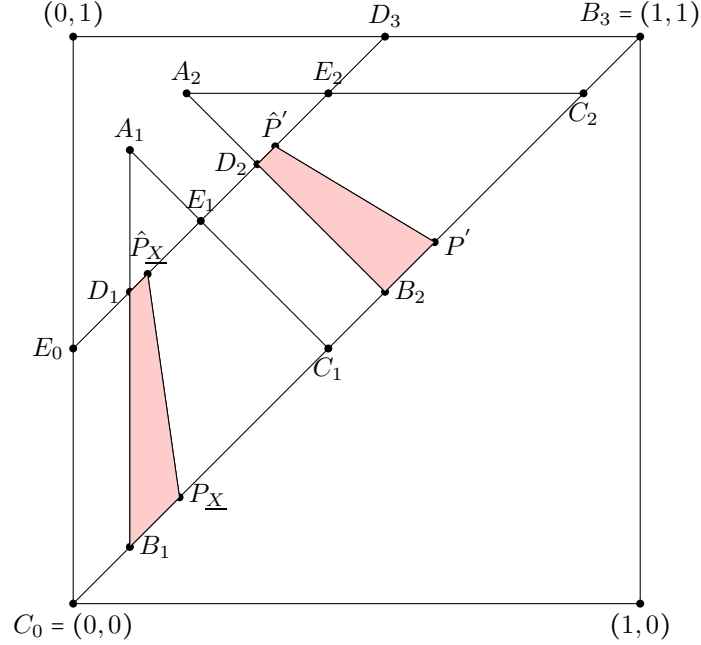


Figure 6: Illustration of (18), where $P_{\underline{X}} \in \overline{B_i C_i}$ for some i . Both red quadrangles have the same measure by Corollary 3 and the intercept theorem.

By Corollary 2,

$$|V_{n,k}(P_{\underline{X}} P_{\overline{Y}} \hat{P}_{\overline{Y}} \hat{P}_{\underline{X}})| \lesssim (c_1(\beta, k) + O(\varepsilon)) n^k$$

and with $\varepsilon \rightarrow 0$ we get (15).

Case b) $r > \frac{q-1}{q^2+q+1}$.

For this case, we use an ordering similar to Case a), but with different polygons, due to the different location of their defining points. Let $A_0 = (0, q(\beta - r))$, $A_{q+1} = ((q+1)r, 1)$ and $I = (0, 1)$. Note that the points $A_i, i = 0, \dots, q+1$, lie on the line $g_{y=x+q(\beta-r)}$. We define the total order \leq as follows:

$$\begin{aligned} V_{n,k,b} = & V_{n,k}(C_0 B_1 A_1 G_1 G_0 A_0 \setminus (\overline{A_1 B_1} \cup \overline{A_1 G_1})) \oplus V_{n,k}(A_1 B_1 C_1 \setminus \overline{A_1 C_1}) \\ & \oplus V_{n,k}(C_1 B_2 A_2 G_2 G_1 A_1 \setminus (\overline{A_2 B_2} \cup \overline{A_2 G_2})) \oplus \dots \\ & \oplus V_{n,k}(C_{q-1} B_q A_q G_q G_{q-1} A_{q-1} \setminus (\overline{A_q B_q} \cup \overline{A_q G_q})) \\ & \oplus V_{n,k}(A_q B_q C_q \setminus \overline{A_q C_q}) \oplus V_{n,k}(C_q B_{q+1} G_{q+1} G_q A_q). \end{aligned}$$

We still have to define the ordering of the elements of $V_{n,k}(C_i B_{i+1} A_{i+1} G_{i+1} G_i A_i)$, $i = 0, \dots, q$, and of $V_{n,k}(A_i B_i C_i)$, $i = 1, \dots, q$ (again, we may allow the complete boundary).

If $(\underline{X}, \overline{X}) = (\underline{Y}, \overline{Y})$ we set in all cases $X \leq_i Y$ if $X \leq_{lex} Y$. Thus let $(\underline{X}, \overline{X}) \neq (\underline{Y}, \overline{Y})$.

First we discuss $V_{n,k}(C_i B_{i+1} A_{i+1} G_{i+1} G_i A_i)$. We divide the hexagon $C_i B_{i+1} A_{i+1} G_{i+1} G_i A_i$ into two quadrangles $C_i B_{i+1} A_{i+1} A_i$ and $A_i A_{i+1} G_{i+1} G_i$, define the corresponding orderings for both quadrangles and then explain how they are combined. The definition of the ordering of $V_{n,k}(C_i B_{i+1} A_{i+1} A_i)$ is similar to $V_{n,k}(C_i B_{i+1} D_{i+1} E_i)$ of Case a), i.e., for $X, Y \in V_{n,k}(C_i B_{i+1} A_{i+1} A_i)$ and $(\underline{X}, \overline{X}) \neq (\underline{Y}, \overline{Y})$ we set

$$X \leq_i Y \text{ if } ((1 - i/q)\underline{X} + i\overline{X}/q, -\underline{X} + \overline{X}) \leq_{lex} ((1 - i/q)\underline{Y} + i\overline{Y}/q, -\underline{Y} + \overline{Y}).$$

Concerning $V_{n,k}(A_i A_{i+1} G_{i+1} G_i)$ we work with polar coordinates $\varphi_i(X)$ and $r_i(X)$ of points $\frac{1}{n}(\underline{X}, \overline{X})$, where $X \in V_{n,k}(A_i A_{i+1} G_{i+1} G_i)$, in the coordinate system with origin I and x -axis in the direction of $\overrightarrow{IA_i}$ and arbitrary, but fixed unit length. Similarly to $V_{n,k}(B_i C_i E_i D_i)$ in Case a), the ordering is a lexicographic ordering with respect to the reflected polar coordinates, i.e., for $X, Y \in V_{n,k}(A_i A_{i+1} G_{i+1} G_i)$ and $(\underline{X}, \overline{X}) \neq (\underline{Y}, \overline{Y})$ we set

$$X \leq_i Y \text{ if } (\varphi_i(X), -r_i(X)) \leq_{lex} (\varphi_i(Y), -r_i(Y)).$$

For a point $P \in C_i B_{i+1} A_{i+1} A_i$ let \tilde{P} be the intersection point of the line $g_{y=x+q(\beta-r)}$ with the line through P that is parallel to $C_i A_i$. If, in particular, $P = \frac{1}{n}(\underline{X}, \overline{X})$ with $X \in V_{n,k}(C_i B_{i+1} A_{i+1} A_i)$, then let $\tilde{\varphi}(X)$ be the angular coordinate of \tilde{P} in the coordinate system introduced for $V_{n,k}(A_i A_{i+1} G_{i+1} G_i)$.

The combination of the two orderings is as follows: Let $X \in V_{n,k}(C_i B_{i+1} A_{i+1} A_i)$ and $Y \in V_{n,k}(A_i A_{i+1} G_{i+1} G_i)$ and $(\underline{X}, \overline{X}) \neq (\underline{Y}, \overline{Y})$. We set

$$X \leq_i Y \text{ if } \tilde{\varphi}_i(X) \leq \varphi_i(Y).$$

Finally, we discuss $V_{n,k}(A_i B_i C_i)$. Here the ordering is a lexicographic ordering of the reflected polar coordinates in the coordinate system with origin A_i and x -axis in the direction of $\overrightarrow{A_i B_i}$ and arbitrary, but fixed unit length.

The whole ordering is illustrated in Figure 7.

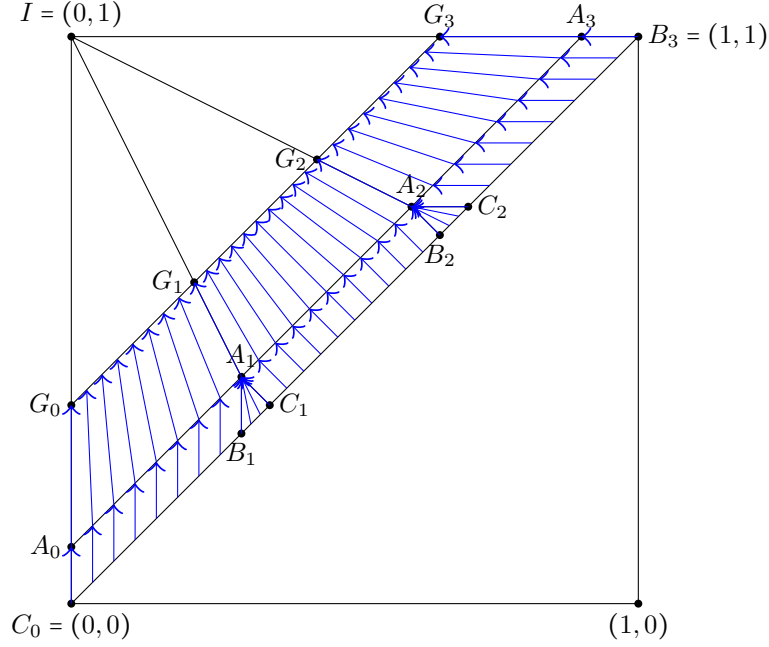


Figure 7: Schematic illustration of the ordering for Case b).

Lemma 15. *Let f be the numbering for Case b). Then*

$$B_f(G) \lesssim (c_2(\beta, k) + c_3(\beta, k))n^k.$$

Proof. Let $P = (\xi, \xi)$ be any point on the segment $\overline{C_0 B_{q+1}}$, i.e., $0 \leq \xi \leq 1$. With P we associate a new point \hat{P} as follows: If $P \in \overline{C_i B_{i+1}}$ for some i , i.e., P belongs to the quadrangle $C_i B_{i+1} A_{i+1} A_i$, then we already defined \hat{P} . The point \hat{P} is the intersection point of the line $g_{y=x+\beta}$ with the line through \hat{P} and I . If $P \in \overline{B_i C_i}$ for some i then let \hat{P} be the intersection point of the line $g_{y=x+\beta}$ with the line through I and A_i , see Figure 8.

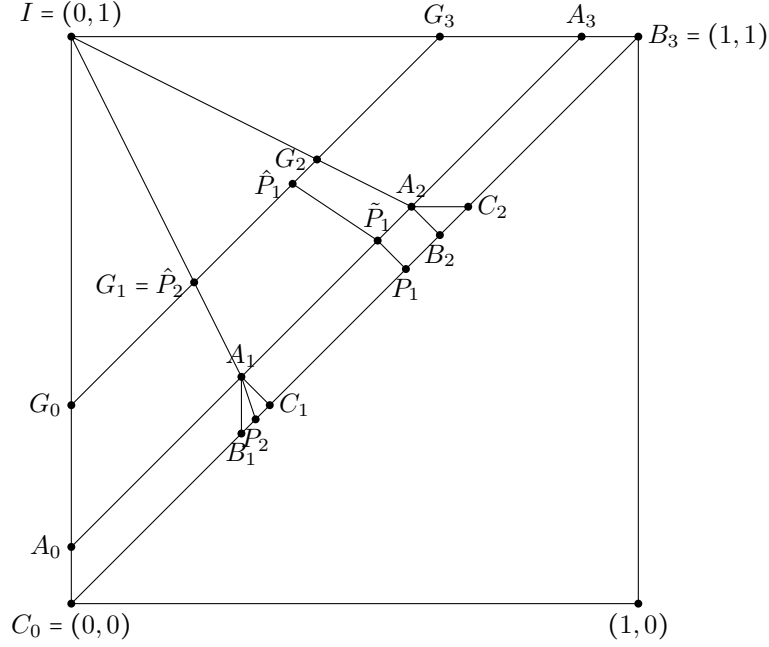


Figure 8: Important points for the upper bound for Case b).

For $X \in V_{n,k,b}$ we define $P_{\underline{X}}$ and $P_{\overline{X}}$ as in (14) and, analogously to Case a), have to prove that for $n \rightarrow \infty$

$$|V_{n,k}(P_{\underline{X}}P_{\overline{Y}}\hat{P}_{\overline{Y}}\hat{P}_{\underline{X}})| \lesssim (c_2(\beta, k) + c_3(\beta, k))n^k. \quad (19)$$

We define P' as in (17). With the same arguments as for Case a) it is sufficient to prove that

$$\mu(P_{\underline{X}}P'\hat{P}'\hat{P}_{\underline{X}}) = c_2(\beta, k) + c_3(\beta, k).$$

Using Lemma 4, one can verify that for all possible i ,

$$\begin{aligned} \mu(P_{\underline{X}}P'\hat{P}'\hat{P}_{\underline{X}}) &= \mu(B_i B_{i+1} A_{i+1} G_{i+1} G_i A_i) = \mu(C_i C_{i+1} A_{i+1} G_{i+1} G_i A_i) \\ &= \mu(F_i F_{i+1} G_{i+1} G_i) + \mu(A_{i+1} F_{i+1} C_{i+1}) - \mu(A_i F_i C_i) \\ &= c_2(\beta, k) + c_3(\beta, k). \end{aligned} \quad (20)$$

An illustration of this fact can be found in Figures 9 and 10.

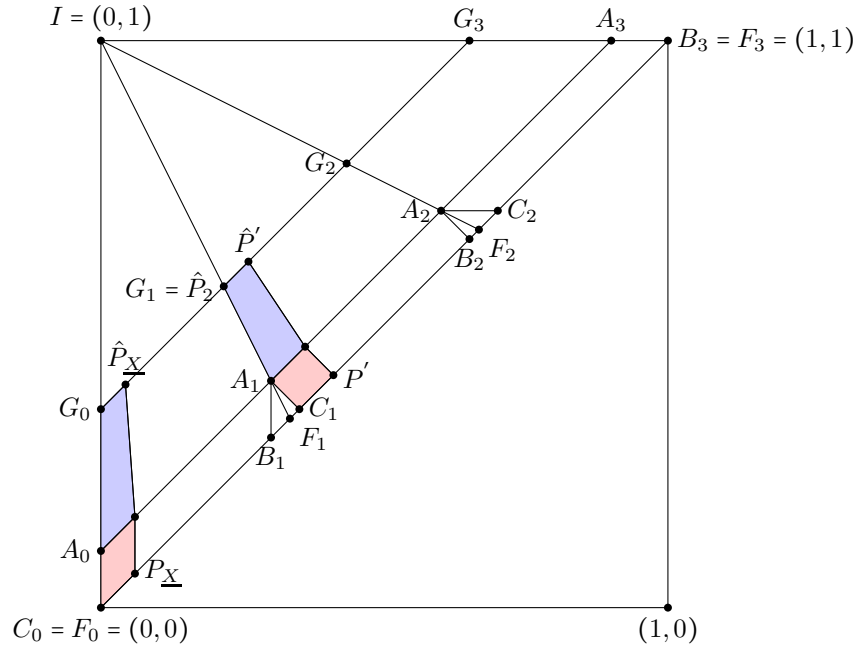


Figure 9: Illustration of (20), where $P_{\underline{X}} \in \overline{C_i B_{i+1}}$ for some i . Both red and both blue quadrangles have the same measure by Corollary 3 and the intercept theorem.

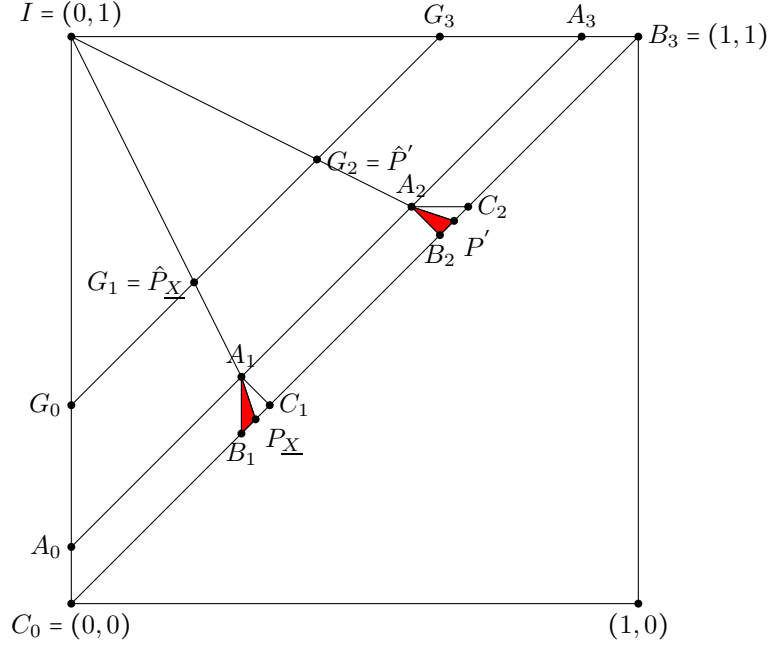


Figure 10: Illustration of (20), where $P_{\underline{X}} \in \overline{B_i C_i}$ for some i . Both red quadrangles have the same measure by Corollary 3.

10 Open problems

We formulate the following conjecture in form of a problem because we are rather convinced that it is correct.

Problem 1. *Prove that the ordering for Case b) presented in Section 9 and illustrated in Figure 7 defines an asymptotically optimal bandwidth numbering.*

A larger program is formulated in the second problem:

Problem 2. *Find and study other interesting graph classes that allow a reduction to the unit square for the asymptotics and lead to interesting and non-trivial orderings on the unit square.*

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